

Web Acceleration Mechanics

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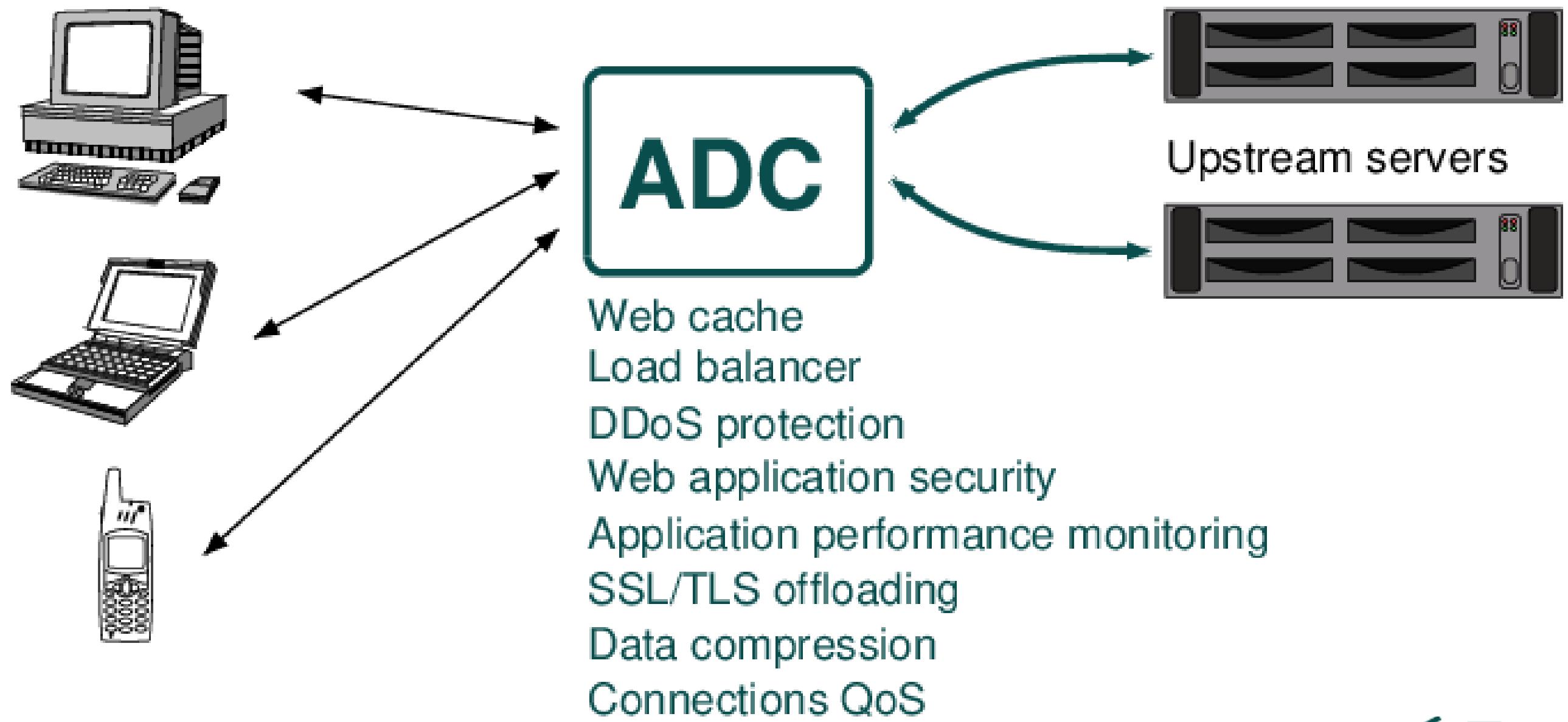
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Who am I?

- ▶ CEO at *Tempesta Technologies, INC*
- ▶ **Custom software development since 2008:**
 - Network security: WAF, VPN, DPI etc.
e.g. *Positive Technologies AF*,
“Visionar” Gartner magic quadrant’15
 - Databases:
one of the top **MariaDB** contributors
 - Performance tuning
- ▶ **Tempesta FW – Linux Application Delivery Controller**

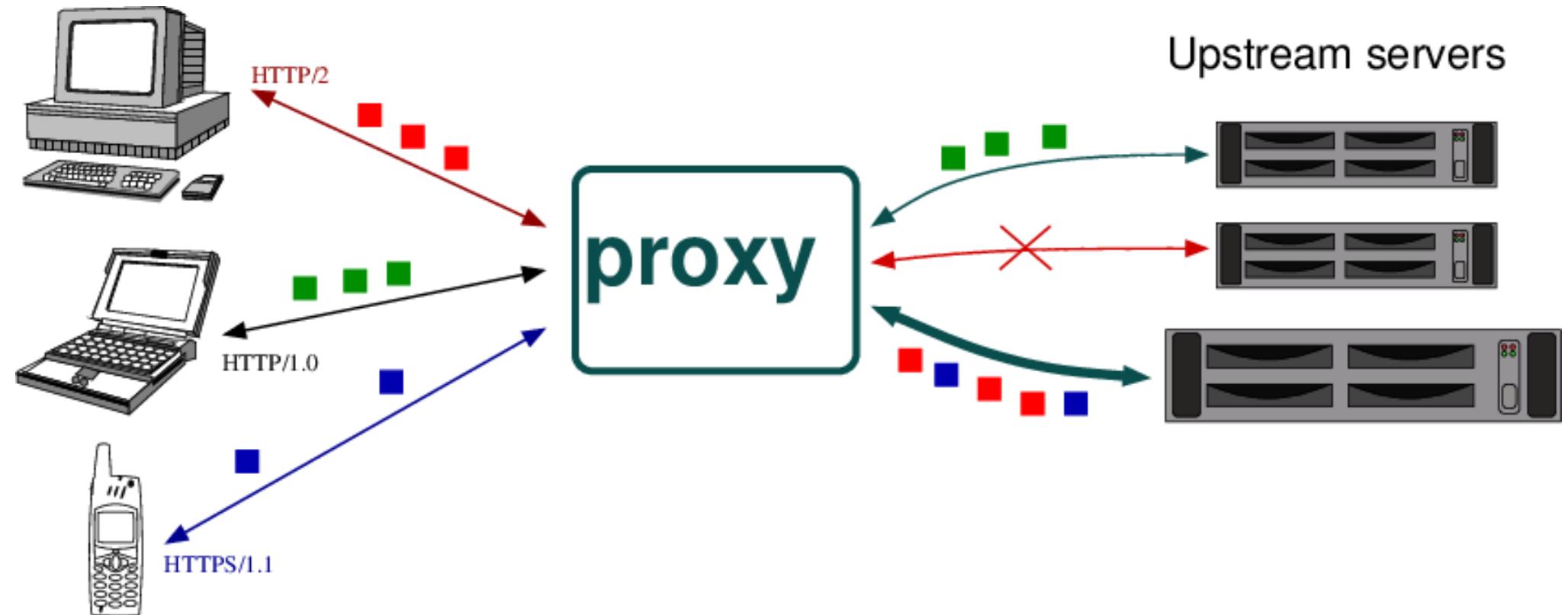


Tempesta FW: Application Delivery Controller (ADC)



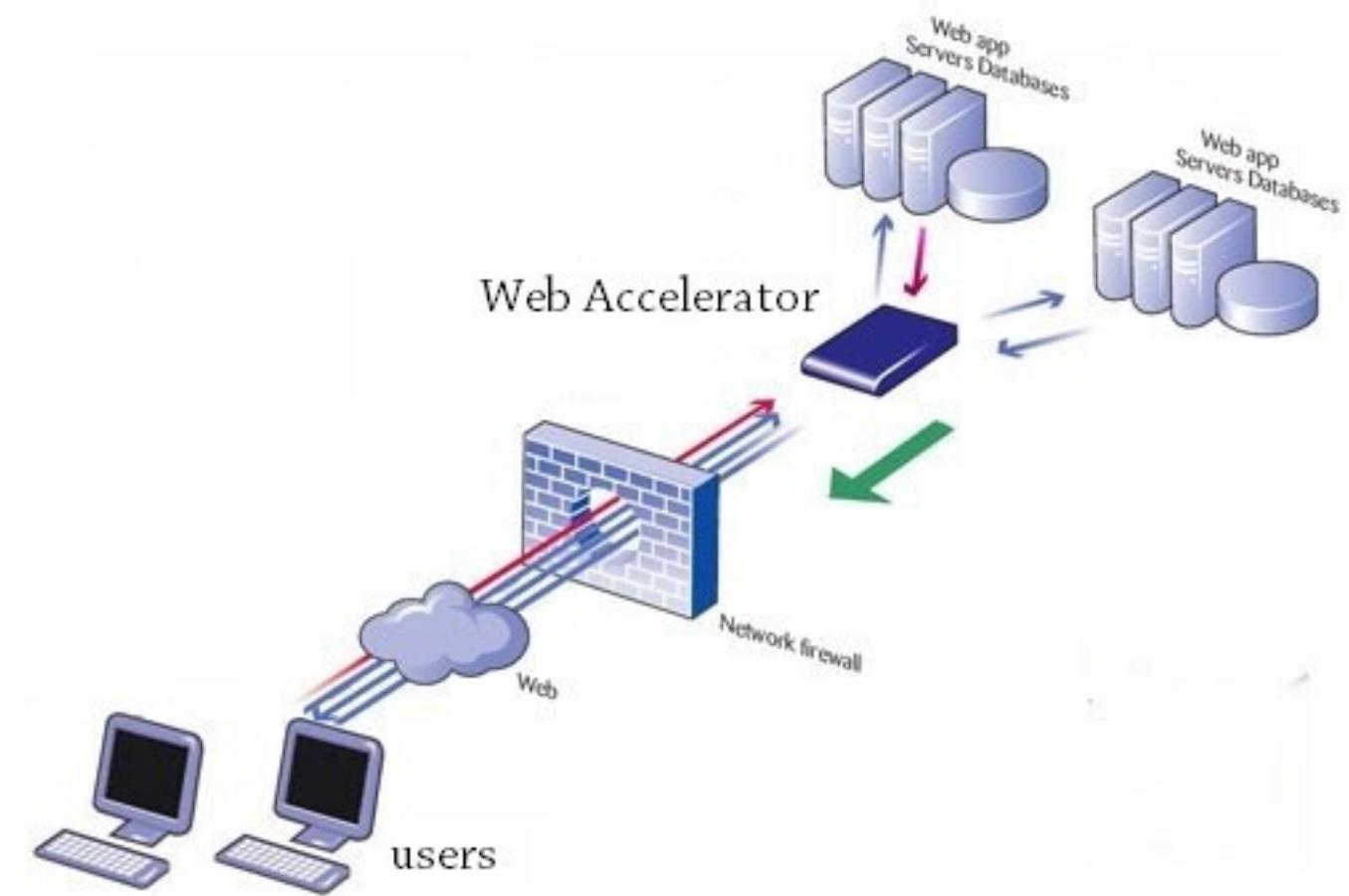
Web accelerator

- ▶ Load balancing
- ▶ Backend connections management
- ▶ TLS termination
- ▶ Protocol downgrade
- ▶ Cache responses



Agenda & Disclaimer

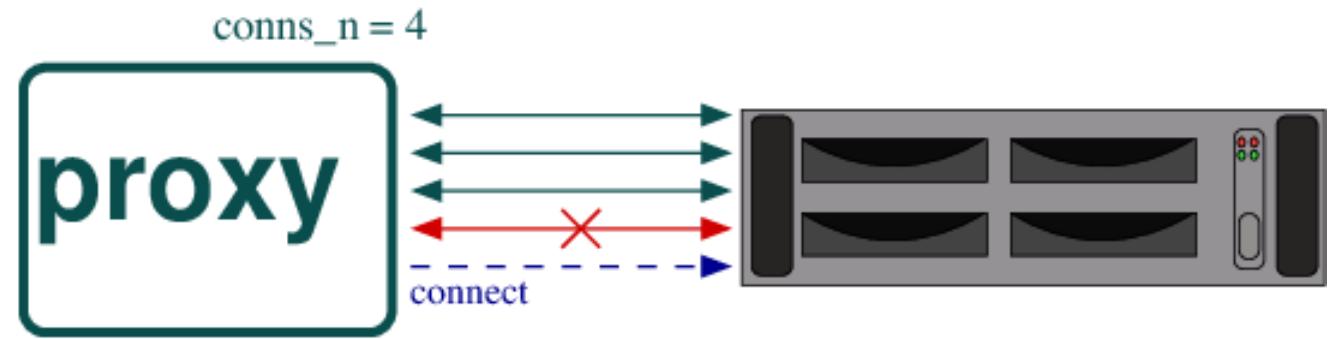
- ▶ Mechanics:
 - **HTTP connections & messages management**
 - **HTTP decoders & parsers**
 - **Web caches**
 - **Network I/O**
 - **Multitasking**
 - **TLS**
- ▶ The web accelerators are mentioned
only as implementation examples
- ▶ Some software is just more familiar to me



HTTP connections & messages management

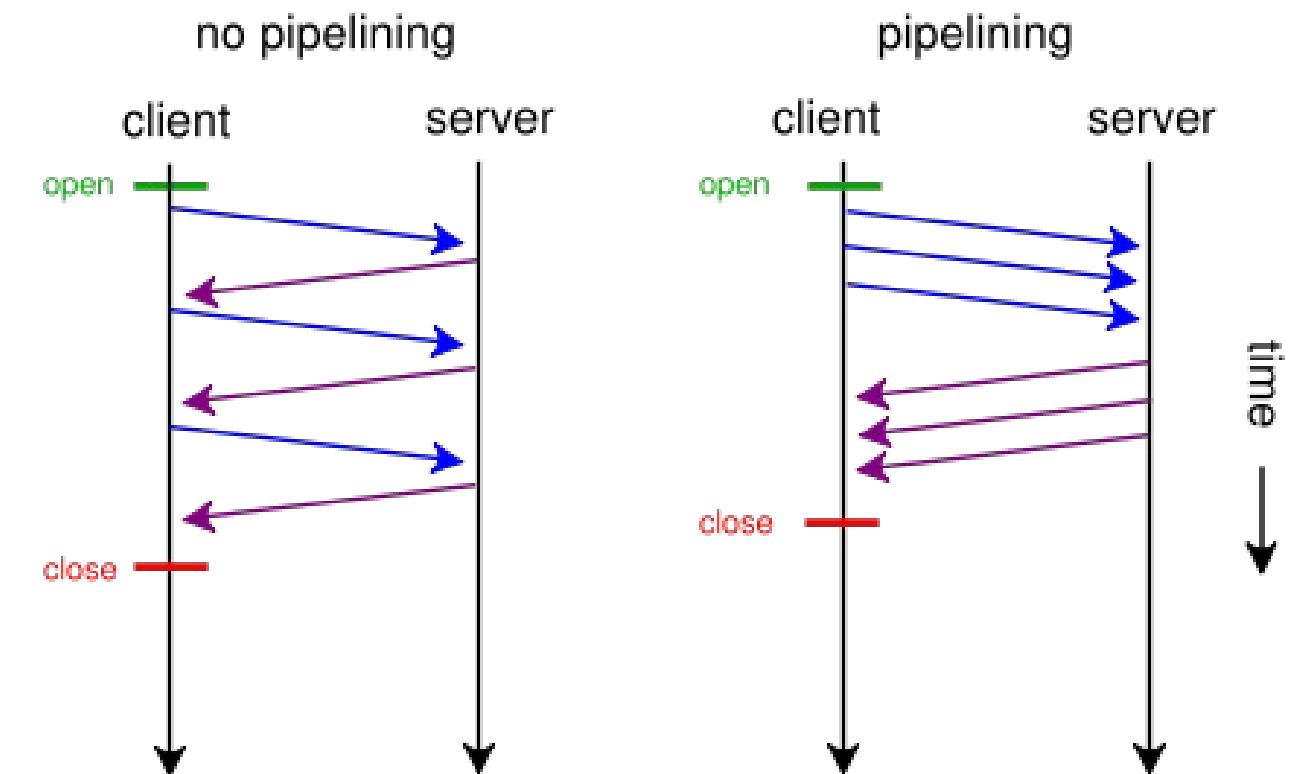
Server connections

- ▶ New connections or **persistent** connections (usual)
- ▶ HTTP keep-alive connections
Keep-Alive: timeout=5, max=10000
- ▶ Reestablish closed KA connection
- ▶ New connections if all are busy
- ▶ **N backend connections = N backend requests in-flight**
- ▶ DDoS and flash crowds:
as many server connections as client connections
- ▶ Run out of port numbers



HTTP/1 message pipelining

- ▶ Mostly unused by proxies
- ▶ Squid, Tempesta FW, Polipo
- ▶ Messages multiplexing
- ▶ Forwarding and reforwarding issues
- ▶ Security issues
 - Breaking authentication
 - HTTP Response splitting



HTTP Response Splitting attack (Web cache poisoning)

[CRLF CRLF]

Client: /redir_lang.jsp?lang=foobar%0d%0aContent-Length:%200%0d%0a%0d%0a
HTTP/1.1%20200%200K%0d%0aContent-Type:%20text/html%0d%0a
Content-Length:%2019%0d%0a%0d%0a<html>Shazam</html>

Client: (one more request to inject Shazam into the cache)

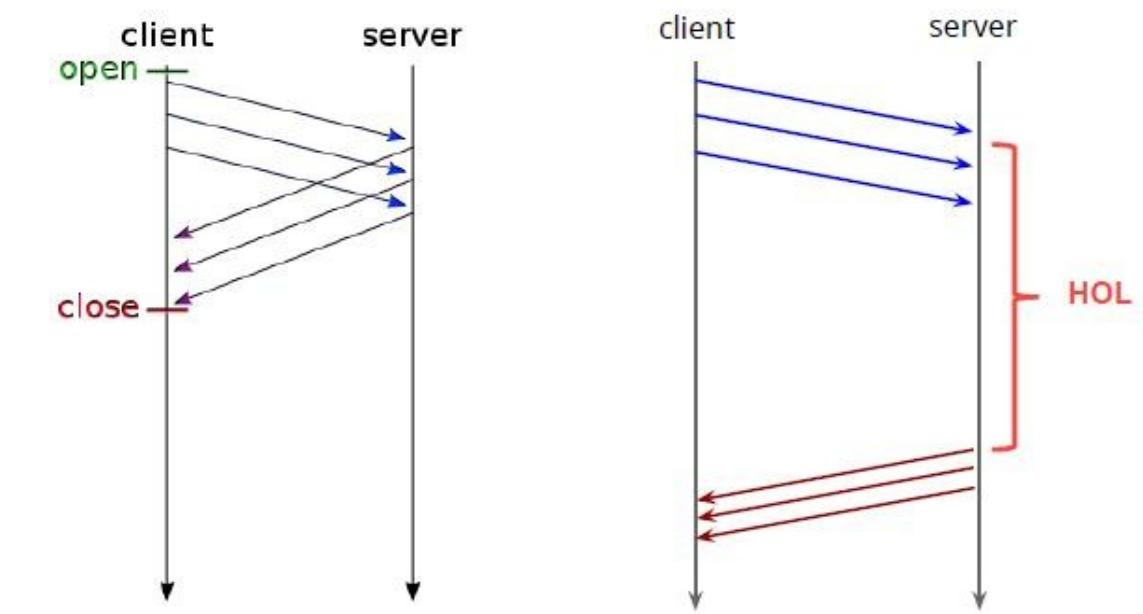
Server: HTTP/1.1 302 Moved Temporarily
Date: Wed, 24 Dec 2003 15:26:41 GMT
Location: http://10.1.1.1/by_lang.jsp?lang=foobar
Content-Length: 0

(server
uses
lang
as is)
HTTP/1.1 200 OK
Content-Type: text/html
Content-Length: 19
<html>Shazam</html>

- ▶ Decode and validate URI (also for injection attacks)
- ▶ **HTTP/2 isn't vulnerable**

HTTP/2

- ▶ Requests multiplexing instead of pipelining
- ▶ **Pros**
 - Responses are sent in *parallel* and in *any order* (no head-of-line blocking)
 - Compression(?)
- ▶ **Cons**
 - Less zero copy (different dynamic tables)
- ▶ HTTP/2 backend connections (H2O, HAProxy, Envoy)
 - Slow backend connections (e.g. CDN)
 - Slow logic (e.g. dynamic content w/ database access)



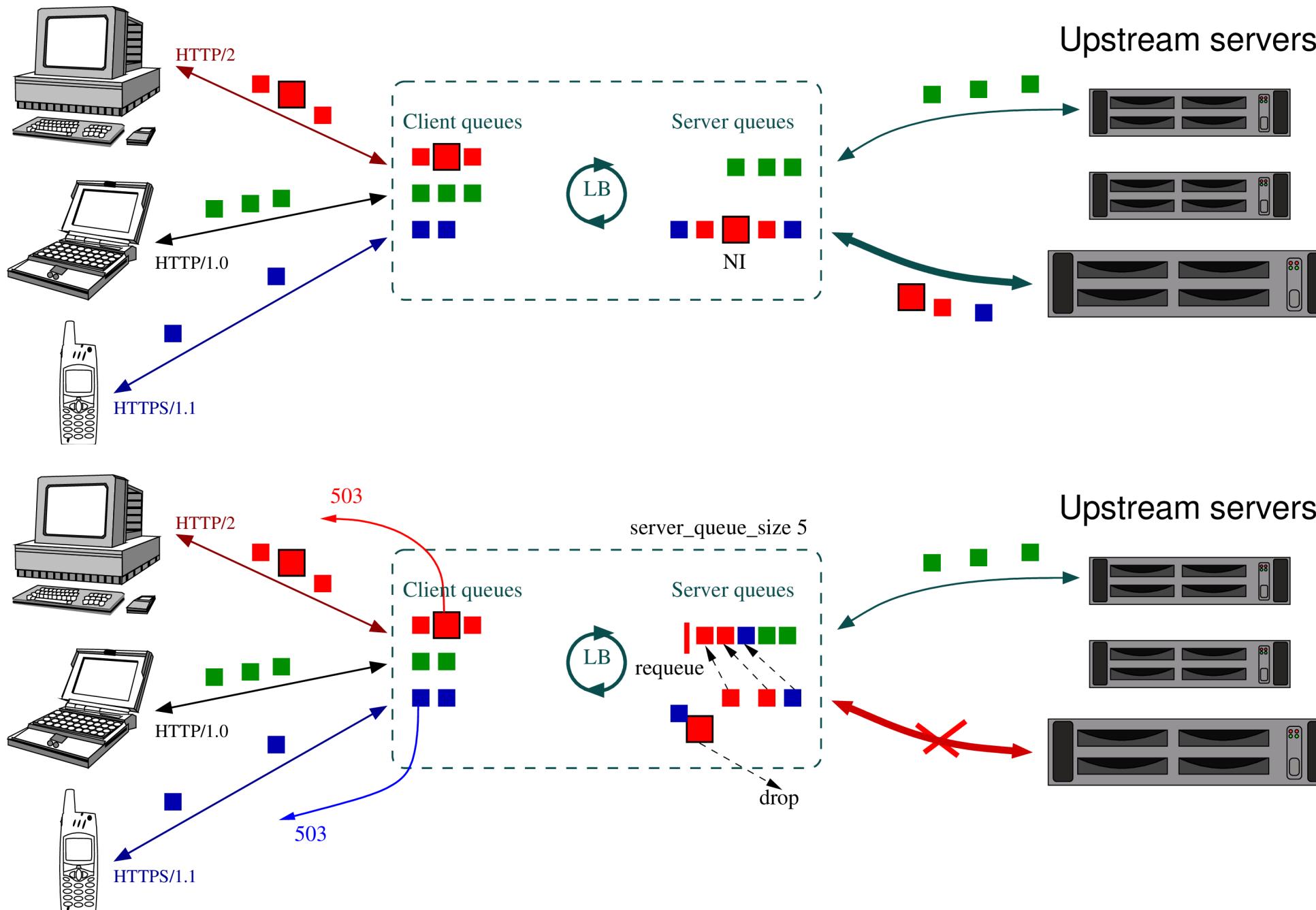
(Non-)indempotent requests

- ▶ RFC 7231 4.2.2: **Non-idempotent** request changes server state
- ▶ Idempotent (safe) methods: **GET, HEAD, TRACE, OPTIONS**
 - User responsibility: GET /action?do=delete
 - Can be POST (e.g. web-search)
- ▶ **Let a user** to specify idempotence of methods and resources:
nonidempotent GET prefix “/action”
- ▶ Cannot be retried in HTTP/1
- ▶ HTTP/2 (RFC 7540 8.1.4): **can retry non-idempotent** requests
 - GOAWAY and RST_STREAM with REFUSED_STREAM

Resend HTTP requests?

- ▶ **Tainted requests:** `server_forward_retries` and `server_forward_timeout`
- ▶ **Request-killers** – RFC 7230 6.3.2: “*a client MUST NOT pipeline immediately after connection establishment*”
 - Connection with re-forwarded requests is **non-schedulable**
- ▶ **Non-idempotent requests** can be reforwarded
 - `server_retry_nonidempotent`
- ▶ Error messages keep **order of responses**

Requests reforwarding (not for persistent sessions!)



Proxy response buffering vs streaming

- ▶ **Buffering**
 - Seems everyone by default
 - Performance degradation on large messages
 - 200 means OK, no incomplete response
- ▶ **Streaming**
 - Varnish, Tengine (patched Nginx)
`proxy_request_buffering & fastcgi_request_buffering`
 - More performance, but 200 doesn't mean full response

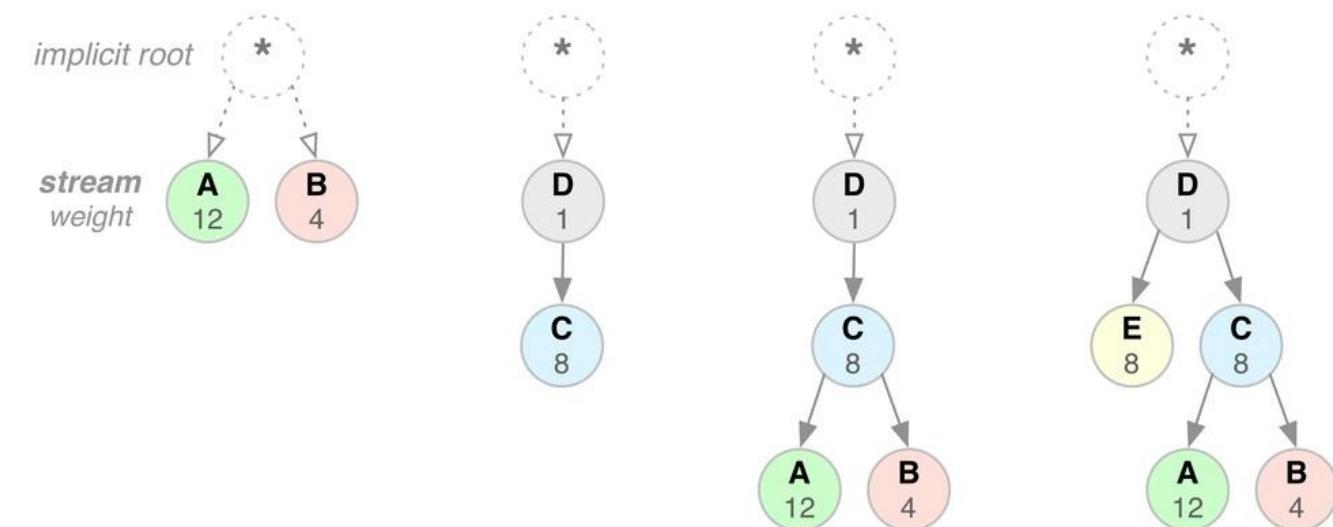
HTTP/2 prioritization

► E.g. A=12, B=4

- Naive: A(12), A(11), A(10), A(9), A(8), A(7), A(6), A(5), A(4), B(4), A(3), B(3),...
- WFQ: A(12), A(11), A(10), B(4), A(9), A(8),...

► Weighted Fair Queue (WFQ)

- H2O approximate in $O(1)$)
- nghttp2 (and Apache) in $O(\log(n))$



Source: High Performance Browser Networking by Ilya Grigorik

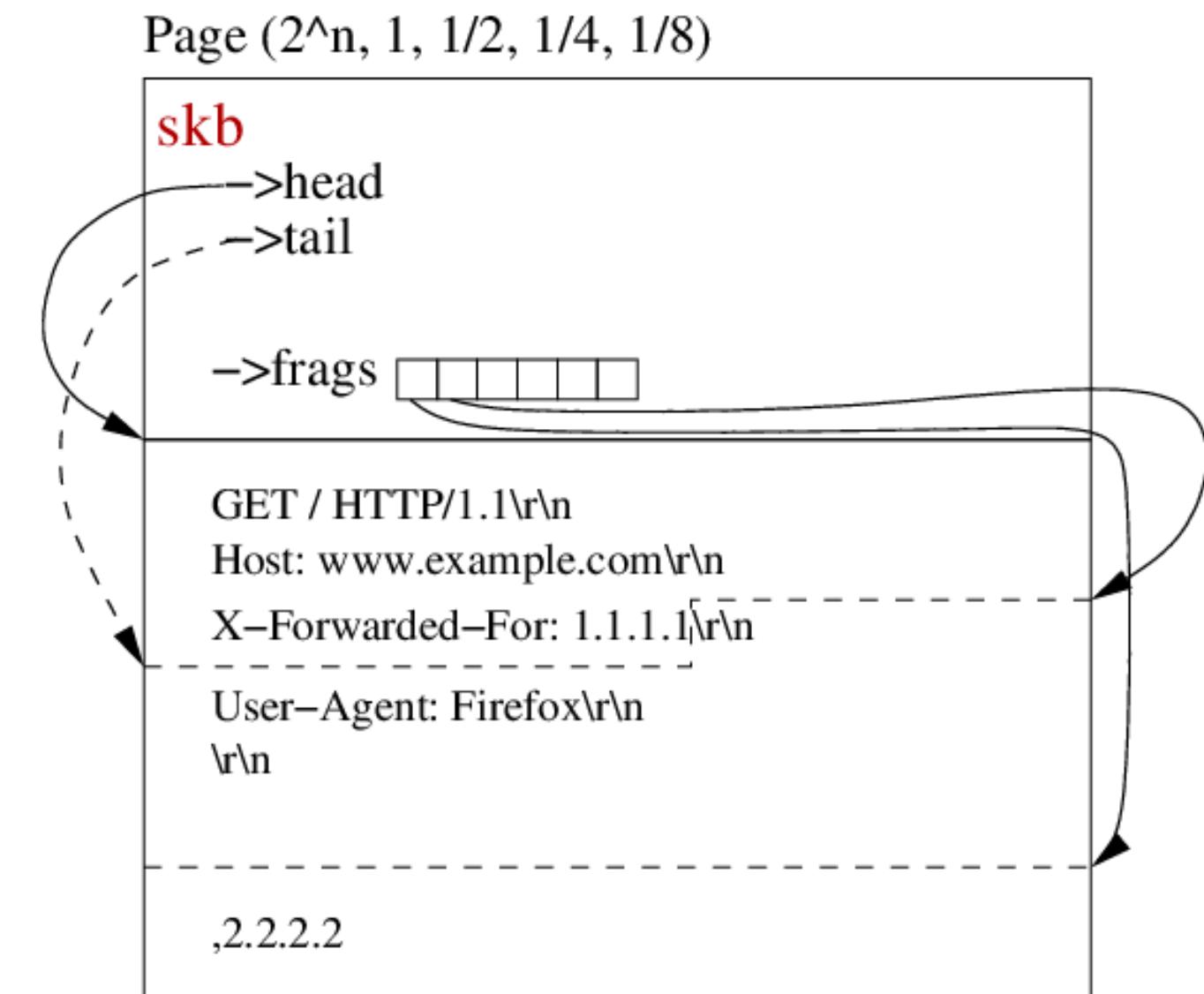
HTTP messages adjustment

- ▶ Add/remove/update HTTP headers
- ▶ Full HTTP message rewriting

GET / HTTP/1.1\r\n
"2.2.2.2" - - - - -
Host: www.example.com\r\nX-Forwarded-For: 1.1.1.1\r\nUser-Agent: Firefox\r\n\r\n

HTTP messages adjustment: zero-copy

- ▶ Tempesta FW: add/remove/update HTTP headers **without copies**
- ▶ Small head and tail memory overheads to avoid dynamic allocations



HTTP Parsers

Multi-pass HTTP parser

```
if ( !strcasecmp( hp->hd[HTTP_HDR_URL].b, "http://", 7) )  
    b = hp->hd[HTTP_HDR_URL].b + 7;  
else if (FEATURE(FEATURE_HTTPS_SCHEME) &&  
    !strcasecmp( hp->hd[HTTP_HDR_URL].b, "https://", 8) )  
b = hp->hd[HTTP_HDR_URL].b + 8;
```

Switch-driven HTTP parser

```
Start: state = 1, *str_ptr = 'b'

    while (++str_ptr) {
        switch (state) { <= check state
            case 1:
                switch (*str_ptr) {
                    case 'a':
                        ...
                        state = 1
                    case 'b':
                        ...
                        state = 2
                }
            case 2:
                ...
        }
    }
}
```

Switch-driven HTTP parser

```
Start: state = 1, *str_ptr = 'b'

    while (++str_ptr) {
        switch (state) {
            case 1:
                switch (*str_ptr) {
                    case 'a':
                        ...
                        state = 1
                    case 'b':
                        ...
                        state = 2 <= set state
                }
            case 2:
                ...
        }
    ...
}
```

Switch-driven HTTP parser

```
Start: state = 1, *str_ptr = 'b'

    while (++str_ptr) {
        switch (state) {
            case 1:
                switch (*str_ptr) {
                    case 'a':
                        ...
                        state = 1
                    case 'b':
                        ...
                        state = 2
                }
            case 2:
                ...
        }
        ... <= jump to while
    }
```

Switch-driven HTTP parser

```
Start: state = 1, *str_ptr = 'b'

    while (++str_ptr) {
        switch (state) { <= check state
            case 1:
                switch (*str_ptr) {
                    case 'a':
                        ...
                        state = 1
                    case 'b':
                        ...
                        state = 2
                }
            case 2:
                ...
        }
    }
}
```

Switch-driven HTTP parser

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Start: state = 1, *str_ptr = 'b'

    while (++str_ptr) {
        switch (state) {
            case 1:
                switch (*str_ptr) {
                    case 'a':
                        ...
                        state = 1
                    case 'b':
                        ...
                        state = 2
                }
            case 2:
                ...      <= do something
        }
    ...
}
```

Switch-driven HTTP parser

```
while (++str_ptr) {  
    switch (state) {  
        case 1:  
            switch (*str_ptr) {  
                case 'a':  
                    ...  
                    state = 1  
                case 'b':  
                    ...  
                    state = 2  
            }  
        case 2:  
            ...  
    }  
    ...  
}
```

The diagram illustrates the state transitions of a switch-driven parser. It starts with a red vertical line labeled '3' at the top, representing the initial state. An arrow points from '3' to a red box labeled '1'. Inside box '1', there is a red box labeled '4'. From '4', a red arrow points down to a red box labeled '2'. From '2', a red arrow points back up to '1'. This indicates a loop between states 1 and 2 based on the character at the string pointer. The code snippet shows that after each iteration, the string pointer is incremented and the state is checked again.

```
while (1):  
    STATE_1:  
        switch (*str_ptr) {  
            case 'a':  
                ...  
                ++str_ptr  
                goto STATE_1  
            case 'b':  
                ...  
                ++str_ptr  
        }  
    STATE_2:  
        ...
```

The diagram illustrates a state transition graph for a parser. It starts with a red arrow pointing down to a red box labeled 'STATE_1'. Inside 'STATE_1', there is a red box labeled '4'. From '4', a red arrow points down to a red box labeled '2'. From '2', a red arrow points back up to '1'. This indicates a loop between STATE_1 and STATE_2 based on the character at the string pointer. The code snippet shows that after each iteration, the string pointer is incremented and the state is checked again. The state transitions are labeled with 'STATE_1' and 'STATE_2'.

HTTP parsers in the wild

- ▶ **Case/switch**
 - Nginx, ATS
- ▶ **Multi-pass (glibc calls)**
 - Varnish, HAProxy
- ▶ **SIMD**
 - Tempesta FW, H2O, CloudFlare
 - Faster on large data (URI, Cookie)
 - Security restrictions against injection attacks (Tempesta FW)
SCALE 17x: <https://www.slideshare.net/AlexanderKrizhanovsky1/fast-http-string-processing-algorithms>

Why HTTP strings matter?

- ▶ Usual URI – just a hotel query

```
https://www.booking.com/searchresults.en-us.html?
aid=304142&label=gen173nr-
1FCAEoggI46AdIM1gEaIkCiAEBmAExuAEZyAEP2AEB6AEB-
AECiAIBqAIDuAKAg4DkBcACAO&sid=686a0975e8124342396dbc1b331
fab24&tmpl=searchresults&ac_click_type=b&ac_position=0&ch
eckin_month=3&checkin_monthday=7&checkin_year=2019&checko
ut_month=3&checkout_monthday=10&checkout_year=2019&class_
interval=1&dest_id=20015107&dest_type=city&dtdisc=0&from_
sf=1&group_adults=1&group_children=0&inac=0&index_postcar
d=0&label_click=undef&no_rooms=1&postcard=0&raw_dest_type
=city&room1=A&sb_price_type=total&sb_travel_purpose=busin
ess&search_selected=1&shw_aparth=1&slp_r_match=0&src=inde
x&srvvid=e0267a2be8ef0020&ss=Pasadena%2C%20California%2C
%20USA&ss_all=0&ss_raw=pasadena&ssb=empty&sshis=0&nflt=hot
elfacility%3D107%3Bmealplan%3D1%3Bpri%3D4%3Bpri
%3D3%3Bclass%3D4%3Bclass%3D5%3Bpopular_activities
%3D55%3Bhr_24%3D8%3Btdb%3D3%3Breview_score
%3D70%3Broomfacility%3D75%3B&rsf=
```

- ▶ How about tons of such queries?
(DDoS)
- ▶ How about injections?

```
/redir_lang.jsp?lang=foobar%0d%0aContent-Length:%200%0d
%0a%0d%0aHTTP/1.1%20200%20OK%0d%0aContent-Type:%20text/
html%0d%0aContent-Length:%2019%0d%0a%0d%0a<html>Shazam</
html>
```

```
case sw_check_uri:
if (usual[ch >> 5] & (1U << (ch & 0x1f)))
    break;
switch (ch) {
case '/':
    r->uri_ext = NULL;
    state = sw_after_slash_in_uri;
    break;
case '.':
    r->uri_ext = p + 1;
    break;
case ' ':
    r->uri_end = p;
    state = sw_check_uri_http_09;
    break;
case CR:
    r->uri_end = p;
    r->http_minor = 9;
    state = sw_almost_done;
    break;
case LF:
    r->uri_end = p;
    r->http_minor = 9;
    goto done;
case '%':
    r->quoted_uri = 1;
...
}
```

Let's check

▶ Reasonable HTTP request

```
./wrk -t 4 -c 128 -d 60s --header 'Connection: keep-alive' --header 'Upgrade-Insecure-Requests: 1'  
--header 'User-Agent: Mozilla/5.0 (X11; Linux x86_64) AppleWebKit/537.36 (KHTML, like Gecko)  
Chrome/52.0.2743.116 Safari/537.36' --header 'Accept: text/html,application/xhtml+xml,  
application/xml;q=0.9,image/webp,*/*;q=0.8' --header 'Accept-Encoding: gzip, deflate, sdch'  
--header 'Accept-Language: en-US,en;q=0.8,ru;q=0.6' --header 'Cookie: a=sdfasd; sdf=3242u389erfhhs;  
djcnjhe=sdfsdafsdjfb324te1267dd' 'http://192.168.100.4:9090/searchresults.en-us.html?  
aid=304142&label=gen173nr-1FCAEoggI46AdIM1gEaIkCiAEBmAExuAEZyAEP2AEB6AEB-AECiAIBqAIDuAKAg4DkBcACAO  
&sid=686a0975e8124342396dbc1b331fab24&tmpl=searchresults&ac_click_type=b&ac_position=0&checkin_month=3&che  
ckin_monthday=7&checkin_year=2019&checkout_month=3&checkout_monthday=10&checkout_year=2019&class_interval=  
1&dest_id=20015107&dest_type=city&dtdisc=0&from_sf=1&group_adults=1&group_children=0&inac=0&index_postcard  
=0&label_click=undef&no_rooms=1&postcard=0&raw_dest_type=city&room1=A&sb_price_type=total&sb_travel_purpos  
e=business&search_selected=1&shw_aparth=1&s1p_r_match=0&src=index&srvid=e0267a2be8ef0020&ss=Pasadena%2C  
%20California%20USA&ss_all=0&ss_raw=pasadena&ssb=empty&sshis=0&nflt=hotelfacility%3D107%3Bmealplan  
%3D1%3Bpri%3D4%3Bpri%3D3%3Bclass%3D4%3Bclass%3D5%3Bpopular_activities%3D55%3Bhr_24%3D8%3Btdb  
%3D3%3Breview_score%3D70%3Broomfacility%3D75%3B&rsf='
```

▶ Even for simple HTTP parser

8.62%	nginx	[.] ngx_http_parse_request_line
2.52%	nginx	[.] ngx_http_parse_header_line
1.42%	nginx	[.] ngx_malloc
0.90%	[kernel]	[k] copy_user_enhanced_fast_string
0.85%	nginx	[.] ngx_strstrn
0.78%	libc-2.24.so	[.] _int_malloc
0.69%	nginx	[.] ngx_hash_find
0.66%	[kernel]	[k] tcp_recvmsg

HTTP/2,3 Decoders

HPACK (HTTP/2) & QPACK (QUIC)

- ▶ Huffman encoding – skewed bytes => no SIMD
- ▶ Dynamic table: at least 4KB overhead for each connection
 - **HPACK bomb (2016)**: N requests with M dynamic table indexed headers of size S => OOM with $N * M * S$ bytes
- ▶ Makes sense for requests only (~1.4% improvement for responses)
<https://blog.cloudflare.com/hpack-the-silent-killer-feature-of-http-2/>
- ▶ Still have to process strings, even for static headers :(

HTTP/2 as the first class citizen

- ▶ HTTP/2 as an add-on to HTTP/1

```
// Cookie is an indexed header in the static table
static ngx_str_t cookie = ngx_string("cookie");
if (ngx_memcmp(header->name.data, cookie.data, cookie.len) == 0)
```

- ▶ Fast headers lookup using static table indexes (H2O, Tempesta FW)
- ▶ Store web cache entries in HTTP/2-friendly format (Tempesta FW)

Web Caching

To cache

- ▶ Determined by Cache-Control and Pragma
- ▶ static (e.g. video, images, CSS, HTML)
- ▶ *some* dynamic
- ▶ Negative results (e.g. 404)
- ▶ Permanent redirects
- ▶ Incomplete results (206, RFC 7233 Range Requests)
- ▶ Methods: **GET**, POST, whatever
- ▶ GET /script?action=delete – this is your responsibility
(but some servers don't cache URIs w/ arguments)

Not to cache

- ▶ Determined by Cache-Control and Pragma
- ▶ Explicit **no-cache** or **private** directives
- ▶ Responses to Authenticated requests
- ▶ Unsafe methods (RFC 7231 4.2.1)
(safe methods: GET, HEAD, OPTIONS, TRACE)
- ▶ Set-Cookie (?)

Cache Set-Cookie?

- ▶ **Varnish, Nginx, ATS don't cache** responses w/ Set-Cookie by default
- ▶ **Tempesta FW, mod_cache, Squid do cache** such responses
- ▶ **RFC 7234, 8 Security Considerations:**
Note that the Set-Cookie response header field [RFC6265] **does not inhibit caching**; a cacheable response with a Set-Cookie header field can be (and often is) used to satisfy subsequent requests to caches. Servers who wish to control caching of these responses are encouraged to **emit appropriate Cache-Control** response header fields.

Cache POST?

- ▶ Discussion: [*https://www.mnot.net/blog/2012/09/24/caching_POST*](https://www.mnot.net/blog/2012/09/24/caching_POST)
 - RFC 7234 4.4: URI must be invalidated
- ▶ Original eBay case: search API with **too many parameters for GET**
[*https://tech.ebayinc.com/engineering/caching-http-post-requests-and-responses/*](https://tech.ebayinc.com/engineering/caching-http-post-requests-and-responses/)
- ▶ **Idempotent** POST (e.g. web-search) – just like GET
- ▶ **Non-idempotent** POST (e.g. blog comment) – cache response for following GET

Cache entries freshness

- ▶ **RFC 7234:** freshness_lifetime > current_age
- ▶ **Freshness calculation** headers: Last-Modified, Date, Age, Expires
- ▶ **Revalidation:**
 - Conditional requests (RFC 7232, e.g. **If-Modified-Since**)
 - Background activity or on-request job
Nginx: proxy_cache_background_update
- ▶ Sometimes it's OK to return **stale cache entries**:
Nginx: proxy_cache_use_stale

HTTP/2 server PUSH

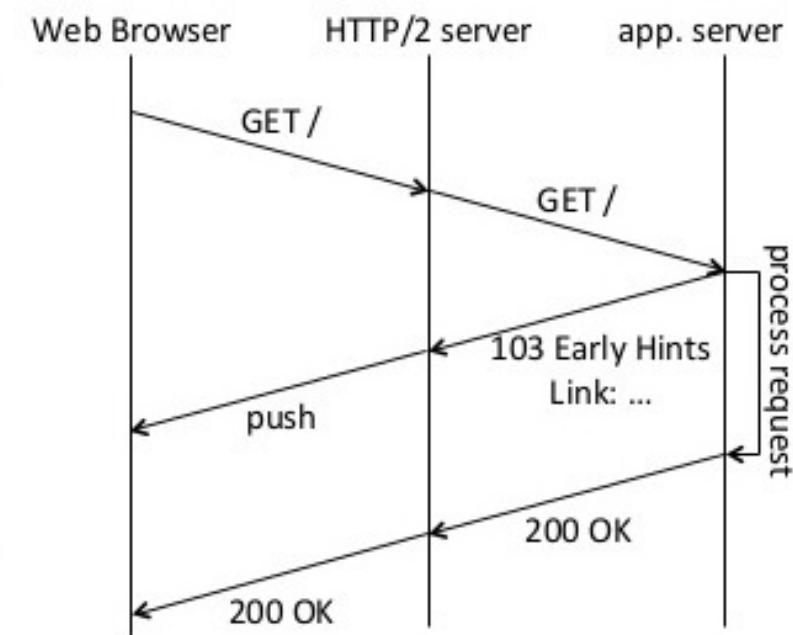
- ▶ **HTTP/2 client ↔ proxy ↔ HTTP/1.1 server** (e.g. H2O, Nginx)
 - Try with `PUSH_PROMISE` and stop by `RST_STREAM`
 - Apache HTTPD: “diary” what has been pushed (`H2PushDiarySize`)
 - H2O CASPer tracks client cache state via cookies
- ▶ **Link header (RFC 5988)**
 - Preload (draft)
<https://w3c.github.io/preload/>
- ▶ **103 Early Hints (RFC 8297, H2O)**

```
GET / HTTP/1.1
Host: example.com

HTTP/1.1 103 Early Hints
Link: </style.css>; rel=preload

HTTP/1.1 200 OK
Content-Type: text/html
Link: </style.css>; rel=preload

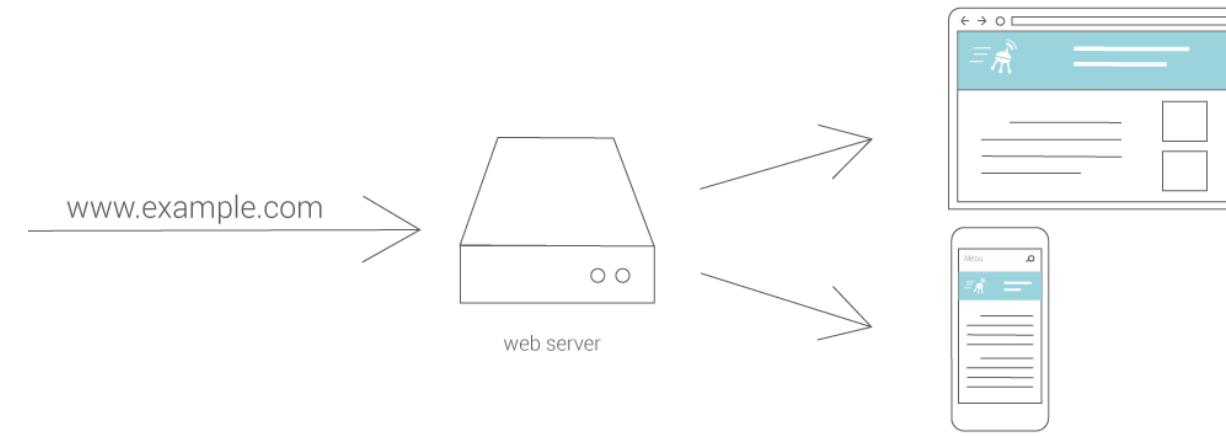
<!DOCTYPE HTML>
...
```



Source: “Reorganizing Website Architecture for HTTP/2 and Beyond” Kazuho Oku

Vary

- ▶ *Accept-Language* – return localized version of page (no need /en/index.html)
- ▶ *User-Agent* – mobile vs desktop (bad!)
- ▶ *Accept-Encoding* – don't send compressed page if browser doesn't understand it
- ▶ **Secondary keys:** say “hello” to databases
- ▶ **Request headers normalization is required!**



Cache storage

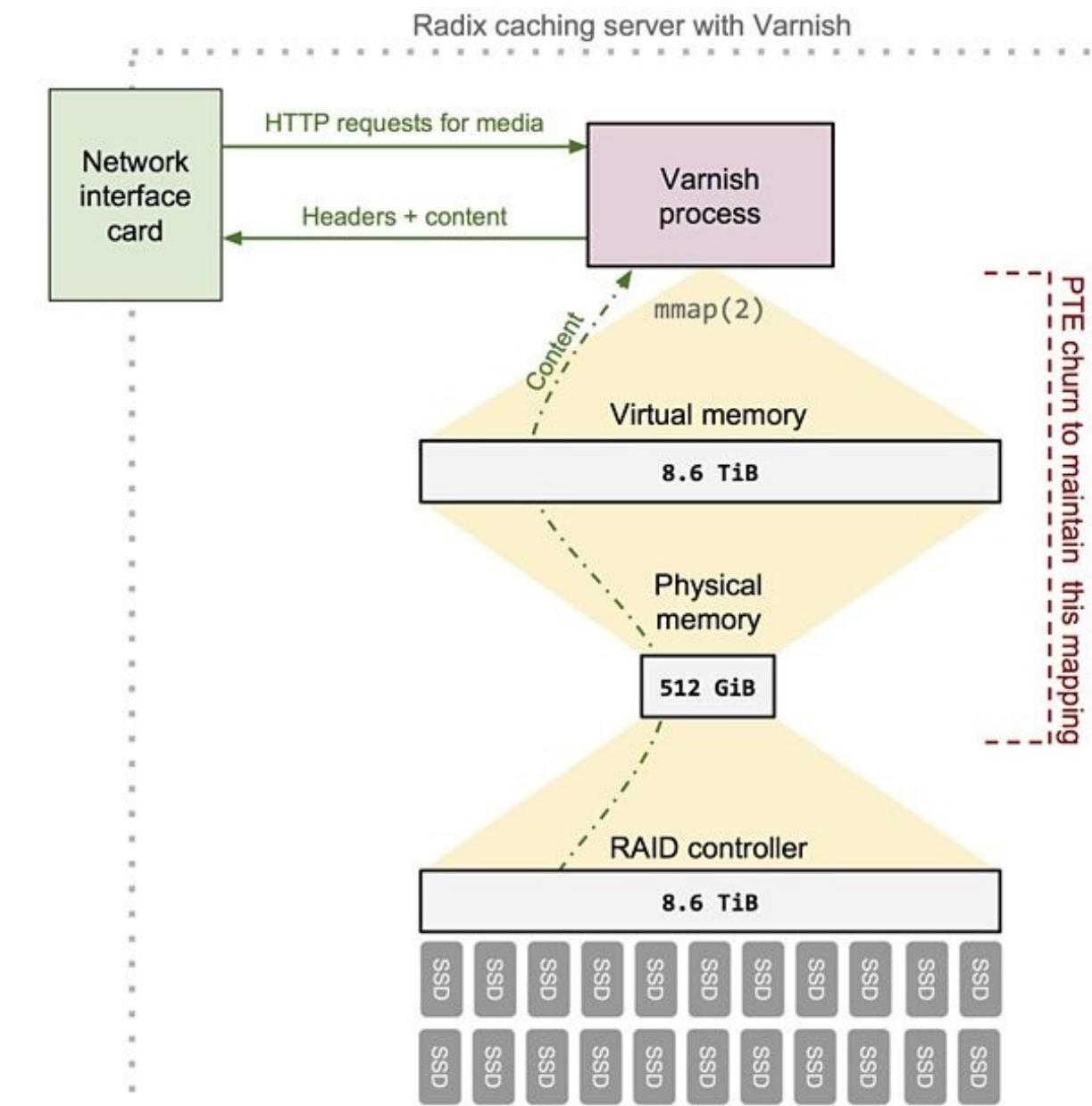
- ▶ **Plain files** (Nginx, Squid, Apache HTTPD)

/cache/0/1d/4af4c50ff6457b8cabfdcd32d0b2f1d0
/cache/5/2e/9f351cdfc8027852656aac5d3f9372e5
/cache/f/22/554a5c654f189c1630e49834c25ae229

- Meta-data in RAM
- Filesystem database
- Easy to manage
- ▶ **Virtual memory** (Varnish): mmap (2) , malloc (3) – no persistency
- ▶ **Database** (Apache Traffic Server, Tempesta FW)
 - Faster access
 - Persistency (no ACID guarantees)

Cache storage: mmap(2)

- ▶ <http://www.bbc.co.uk/blogs/internet/entries/17d22fb8-cea2-49d5-be14-86e7a1dcde04>
- ▶ **48 CPUs, 512GB RAM, 8TB SSD**
- ▶ Basically the same issues with any `mmap()`-based database
<https://www.percona.com/live/data-performance-conference-2016/sessions/linux-kernel-extension-databases>



Requests coalescing

- ▶ Varnish, Nginx
- ▶ Cases:
 - Cold web cache
 - Large responses (e.g. video streaming)
- ▶ Reduces thundering herd of requests to an upstream

Network I/O & Multitasking

IO & multitasking

- ▶ Updated “*Choosing A Proxy Server*”, ApacheCon 2014, Bryan Call

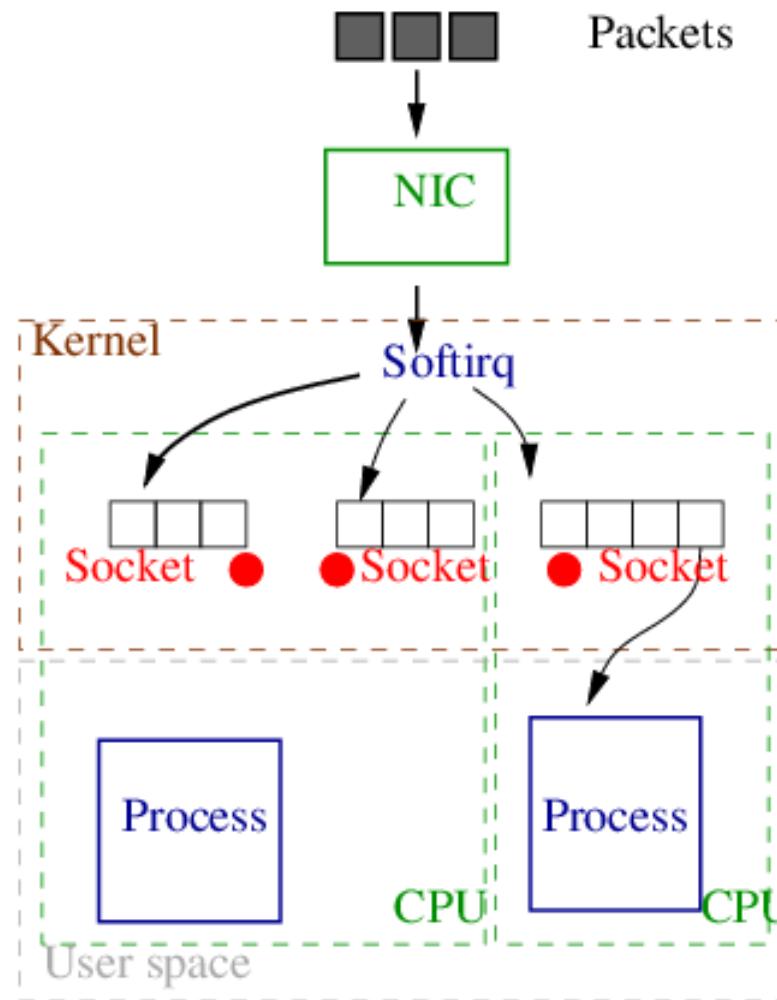
	ATS	Nginx	Squid	Varnish	Apache HTTPD	Tempesta
Threads	X	*		per-session!	X	
Events	X	X	X	partial	X	data-driven
Processes		X	X		X	
Softirq						X

* Nginx 1.7.11 introduced thread pool, but for asynchronous I/O only

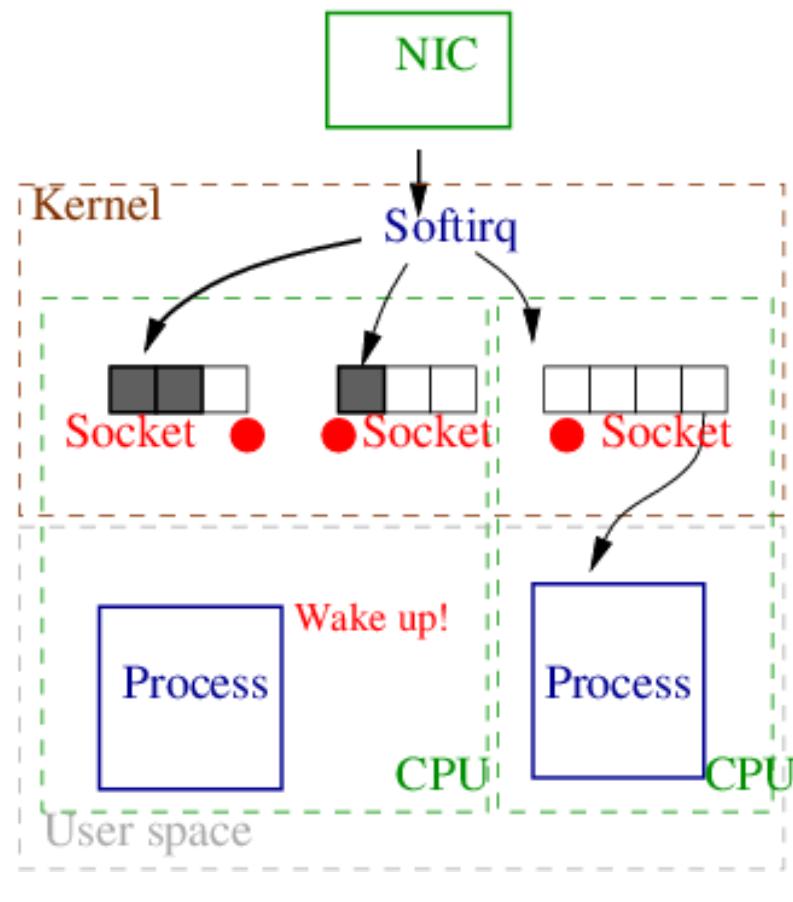
Threads vs Processes vs Softirq

- ▶ Thread per session (e.g. Varnish)
 - Threads overhead: ~6 pages, log(n) scheduler
 - **No slow request starvation** (e.g. asynchronous DDoS)
- ▶ **Meltdown** (reading kernel memory from user space):
CONFIG_PAGE_TABLE_ISOLATION (KPTI, **default** in modern kernels)
 - no lazy TLB as previously, PCID instead
 - **~40%** perf degradation
(<https://mariadb.org/myisam-table-scan-performance-kpti/>)
 - Tempesta FW: no degradation for the in-kernel workload
 - Do you need the protection for a single-user edge server?

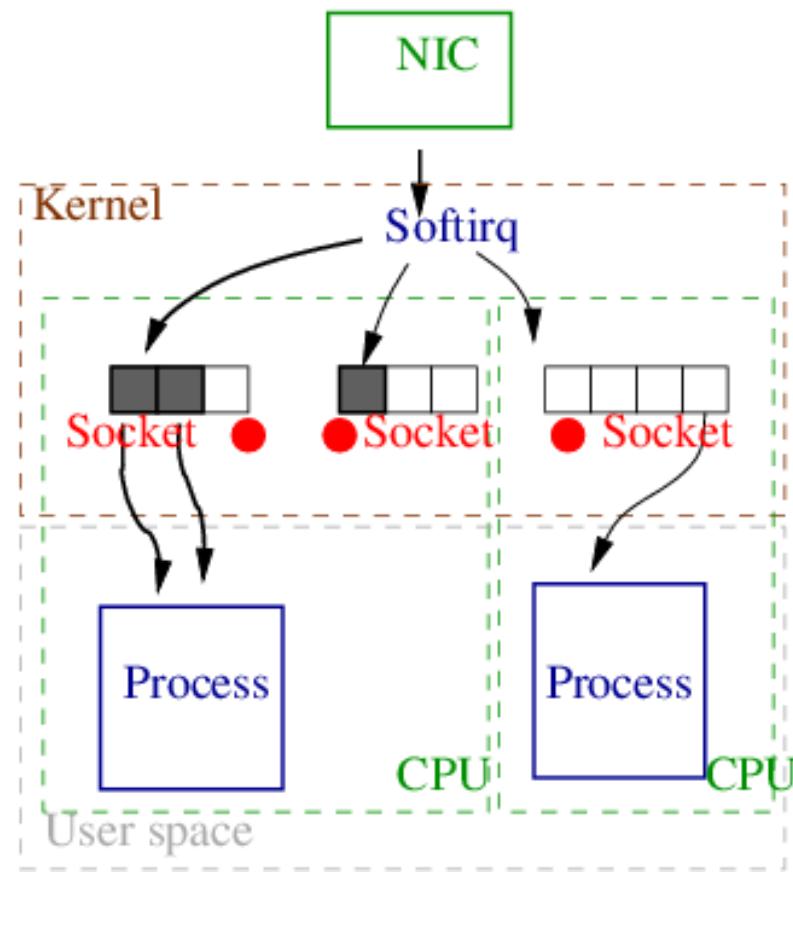
Network asynchronous I/O



Network asynchronous I/O

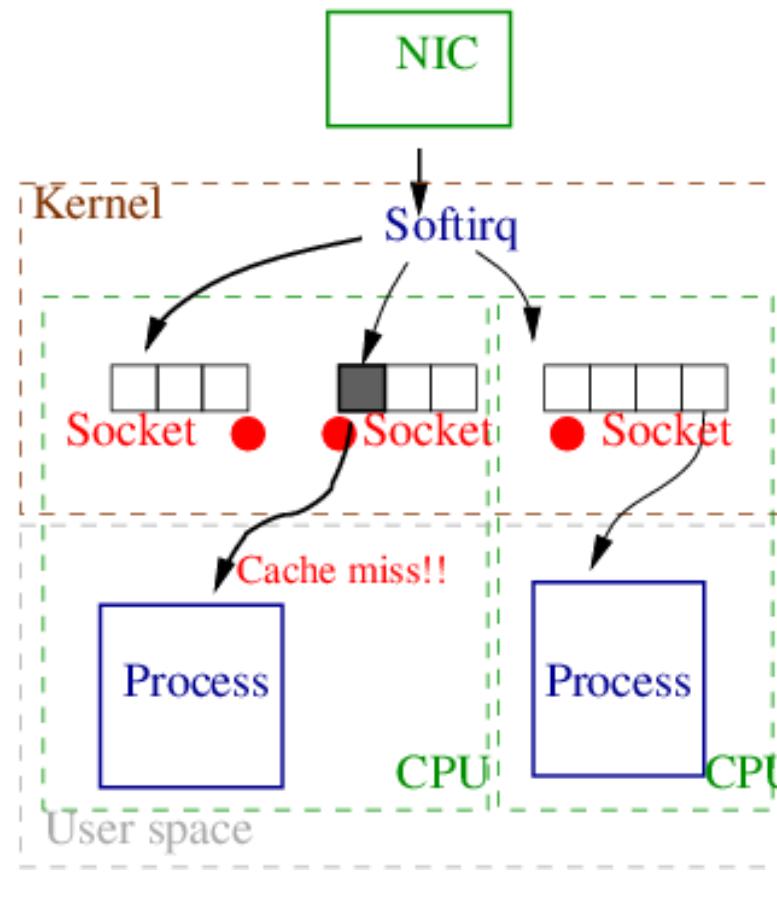


Network asynchronous I/O



Network asynchronous I/O

- ▶ Web cache also resides in CPU caches and **evicts** requests



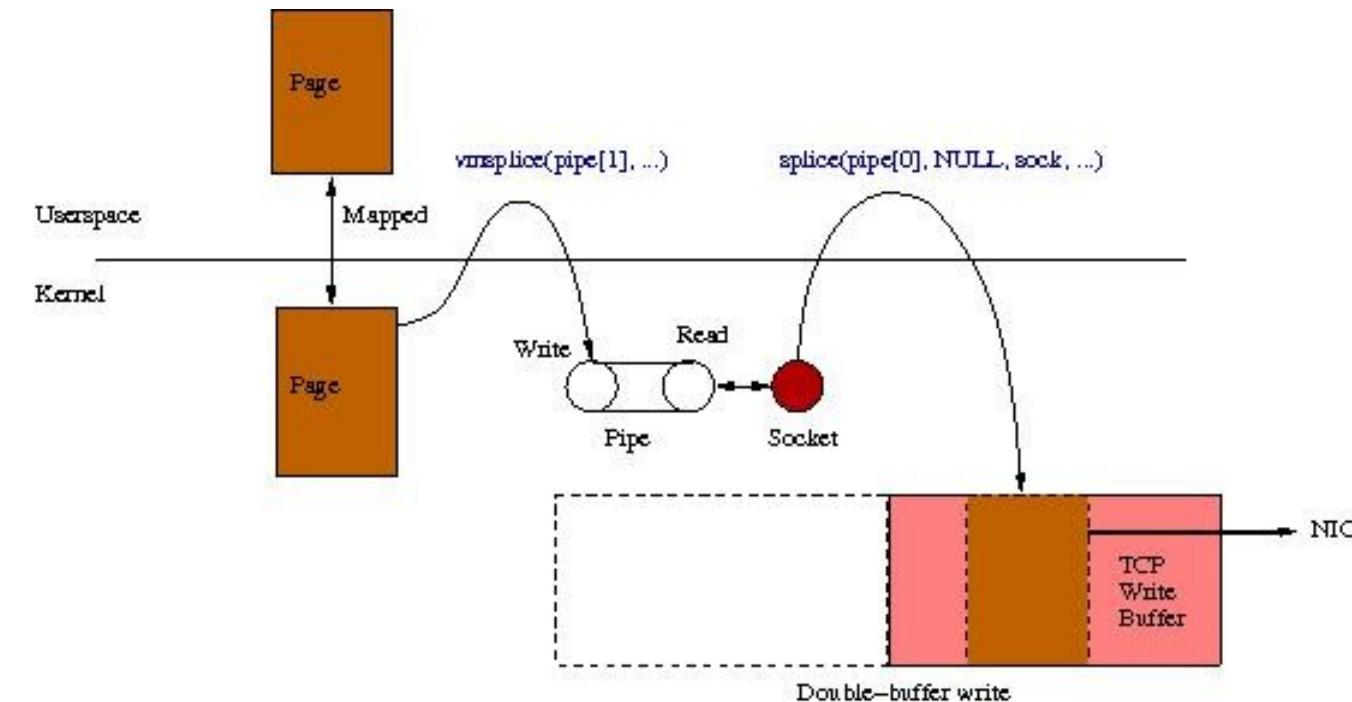
vmsplice(2) zero-copy transmission

- ▶ 2 system calls instead of 1 (**KPTI!!**)
- ▶ Double buffering
- ▶ For large transfers only
- ▶ Examples: HAProxy

```
vmsplice(pipe_[1], iov, num, 0);
for (int r; data_len; data_len -= r)
    r = splice(pipe_[0], NULL, sd, NULL, data_len, 0);
```

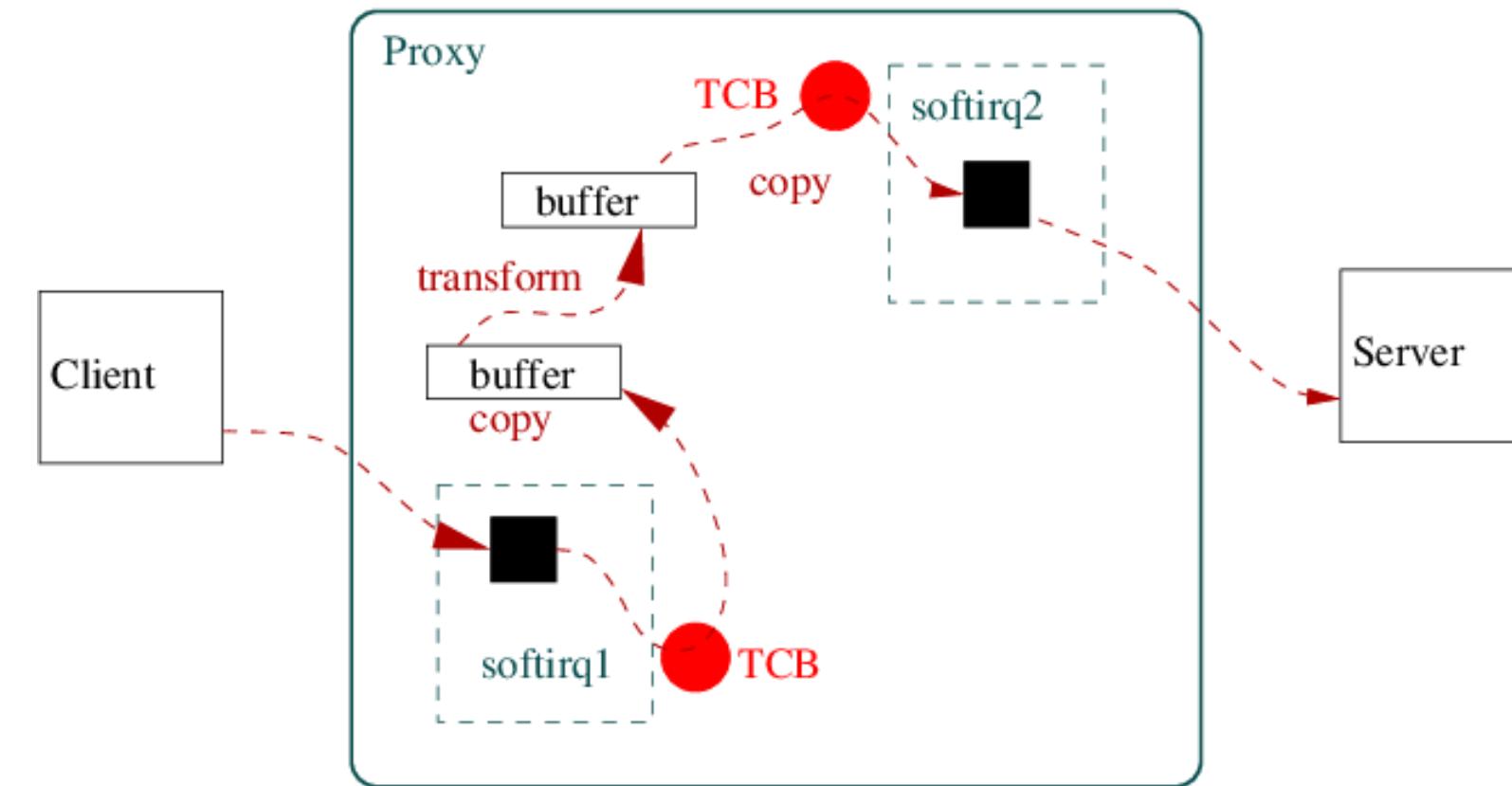
```
# ./nettest/xmit -s65536 -p1000000 127.0.0.1 5500
xmit: msg=64kb, packets=1000000 vmsplice() -> splice()
usr=9259, sys=6864, real=27973
```

```
# ./nettest/xmit -s65536 -p1000000 -n 127.0.0.1 5500
xmit: msg=64kb, packets=1000000 send()
usr=8762, sys=25497, real=34261
```



User space HTTP proxying

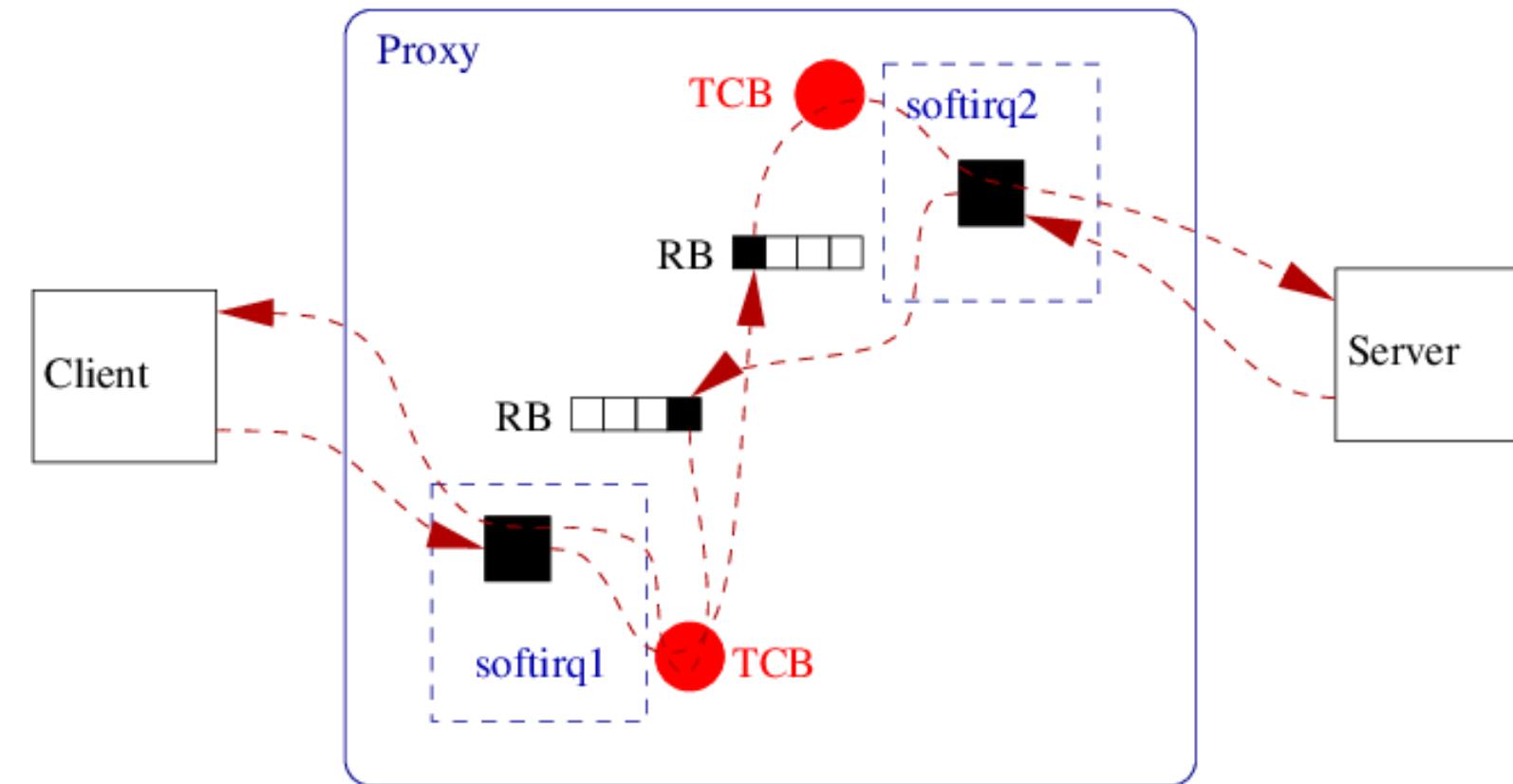
1. Receive request at CPU1
2. Copy request to user space
3. Update headers
4. Copy request to kernel space
(except HAProxy & Tempesta FW)
5. Send the request from CPU2



- ▶ **>= 3 data copies**
- ▶ Access TCP control blocks and data buffers from **different CPUs**

Tempesta FW: zero-copy proxying

- ▶ Socket callbacks call TLS and HTTP processing
- ▶ Everything is processing in softirq (while the data is hot)
- ▶ No receive & accept queues
- ▶ No file descriptors
- ▶ Less locking
- ▶ Lock-free inter-CPU transport
- ▶ => **faster socket reading**
- ▶ => **lower latency**



Logging

- ▶ ATS, Nginx, Squid, Apache HTTPD: *write(2)*
- ▶ Varnish: logs in shared memory → varnishlog
- ▶ Tempesta FW: dmesg (in-memory in further releases)

TLS

TLS termination

- ▶ Native TLS termination: Nginx, HAProxy, H2O, ATS,...
- ▶ Varnish Hitch (no TLS in Varnish cache by design)
- ▶ Tempesta TLS – kernel TLS termination
<https://netdevconf.info/0x14/session.html?talk=performance-study-of-kernel-TLS-handshakes>
- ▶ Intel QuickAssist Technology (QAT) – crypto & compression acceleration
 - PCIe adapters, 89xx chipset, Xeon N
 - OpenSSL & zLib & Nginx patches + user-space library

SSL/TLS: (zero-)copying

- ▶ User-kernel space copying
 - Copy network data to user space
 - Encrypt/decrypt it
 - Copy the data to kernel for transmission
- ▶ **Kernel-mode TLS** (Linux kTLS)
Facebook: <https://lwn.net/Articles/666509/>
 - Eliminates ingress & egress data copyings
 - **Nobody** in user space uses it
(OpenSSL patches are in progress!)
 - Unaware about TCP transmission state (cwnd & rwnd)

TLS record size & TCP

- ▶ TLS record is 16KB by default => several TCP segments
- ▶ TLS decrypts only full records
- ▶ TCP congestion & receive windows can cause last segment delays

```
    ▾ [8 Reassembled TCP Segments (11221 bytes): #169(1460), #170(1460), #172(1460), #174(1460),
        #175(1460), #177(1460), #179(1460), #180(1001)]
        [Frame: 169, payload: 0-1459 (1460 bytes)]
        [Frame: 170, payload: 1460-2919 (1460 bytes)]
        [Frame: 172, payload: 2920-4379 (1460 bytes)]
        [Frame: 174, payload: 4380-5839 (1460 bytes)]
        [Frame: 175, payload: 5840-7299 (1460 bytes)]
        [Frame: 177, payload: 7300-8759 (1460 bytes)]
        [Frame: 179, payload: 8760-10219 (1460 bytes)]
        [Frame: 180, payload: 10220-11220 (1001 bytes)]
        [Segment count: 8]
        [Reassembled TCP length: 11221]
    ▾ Secure Sockets Layer
        ▾ TLSv1 Record Layer: Application Data Protocol: http
            Content Type: Application Data (23)
            Version: TLS 1.0 (0x0301)
            Length: 11216
            Encrypted Application Data: 07ed92e420530da2e2755a5b5372ef32b53e0d4e7c20c3d8...
```

Source: <https://hpbn.co/transport-layer-security-tls/>

TCP & TLS dynamic record size

- ▶ TCP CWND & RWND must be used to avoid multiple RTTs
- ▶ **Typical dynamic TLS record size strategies**
 - Static size (Nginx)
 - Dynamic (HAProxy, ATS, H2O): static algorithms (no cwnd)
 - kTLS: depends on available wmem for the socket
- ▶ QUIC: TLS record = QUIC packet (no need for TLS dynamic records)
- ▶ **Tempesta FW** (pros of being in-TCP/IP stack)
 - TCP send queue, NET_TX_SOFTIRQ
 - min(sndbuff, cwnd, rwnd)
 - http://www.netdevconf.org/2.1/papers/https_tcpip_stack.pdf

References

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- ▶ **HTTP requests proxying**, Alexander Krizhanovsky, <https://natsys-lab.blogspot.com/2018/03/http-requests-proxying.html>
- ▶ **HTTP Strings Processing Using C, SSE4.2 and AVX2**, Alexander Krizhanovsky, <https://natsys-lab.blogspot.com/2016/10/http-strings-processing-using-c-sse42.html>
- ▶ **Fast Finite State Machine for HTTP Parsing**, Alexander Krizhanovsky, <https://natsys-lab.blogspot.com/2014/11/the-fast-finite-state-machine-for-http.html>
- ▶ **Reorganizing Website Architecture for HTTP/2 and Beyond**, Kazuho Oku, <https://www.slideshare.net/kazuho/reorganizing-website-architecture-for-http2-and-beyond>
- ▶ **Server Implementations of HTTP/2 Priority**, Kazuhiko Yamamoto, <https://www.mew.org/~kazu/material/2015-http2-priority2.pdf>
- ▶ **NGINX structural enhancements for HTTP/2 performance**, CloudFlare, <https://blog.cloudflare.com/nginx-structural-enhancements-for-http-2-performance/>

Thanks!

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Custom software development: *tempesta-tech.com/c++-services*

The screenshot shows the GitHub repository page for `tempesta-tech/tempesta`. The top navigation bar includes links for Pull requests, Issues, Marketplace, and Explore. Below the header, there's a summary of repository statistics: 166 issues, 6 pull requests, 0 projects, and 285 forks. The repository is described as a "Linux Application Delivery Controller" with a homepage at <http://tempesta-tech.com/>. A "Manage topics" button is visible at the bottom right of the repository summary.